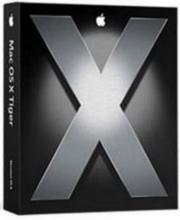
Chapter 10

Operating Systems











Chapter Goals

- Describe the two main responsibilities of an operating system
- Define memory and process management
- Explain how timesharing creates the virtual machine illusion
- Explain the relationship between logical and physical addresses
- Compare and contrast memory management techniques

Software Categories

• Application software Software written to address specific needs—to solve problems in the real world

Word processing programs, games, inventory control systems, automobile diagnostic programs, and missile guidance programs are all application software

• System software Software that manages a computer system at a fundamental level

It provides the tools and an environment in which application software can be created and run

Operating System

- An operating system
 - manages computer resources, such as memory and input/output devices
 - provides an interface through which a human can interact with the computer
 - allows an application program to interact with these other system resources

Operating System

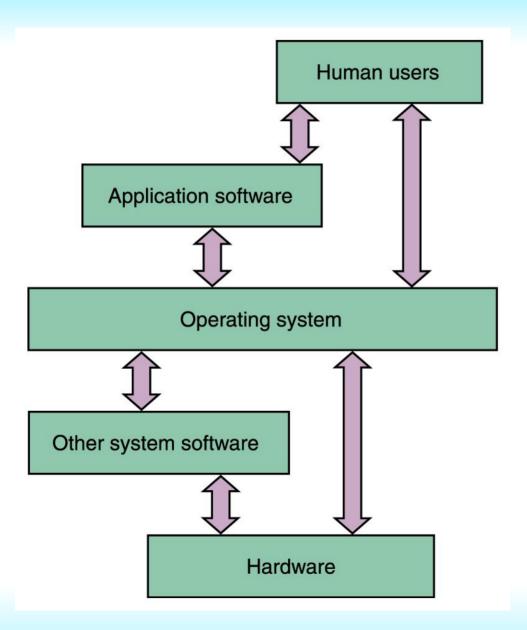


Figure 10.1

An operating system interacts with many aspects of a computer system.

Operating System

- The various roles of an operating system generally revolve around the idea of "sharing nicely"
- An operating system manages resources, and these resources are often shared in one way or another among programs that want to use them

Resource Management

- Multiprogramming The technique of keeping multiple programs in main memory at the same time that compete for access to the CPU so that they can execute
- Memory management The process of keeping track of what programs are in memory and where in memory they reside

Resource Management

- **Process** A program in execution
- The operating system performs **process management** to carefully track the progress of a process and all of its intermediate states
- CPU scheduling determines which process in memory is executed by the CPU at any given point

Batch Processing

- A typical computer in the 1960s and '70s was a large machine
- Its processing was managed by a human operator
- The operator would organize various jobs from multiple users into *batches*

Batch Processing

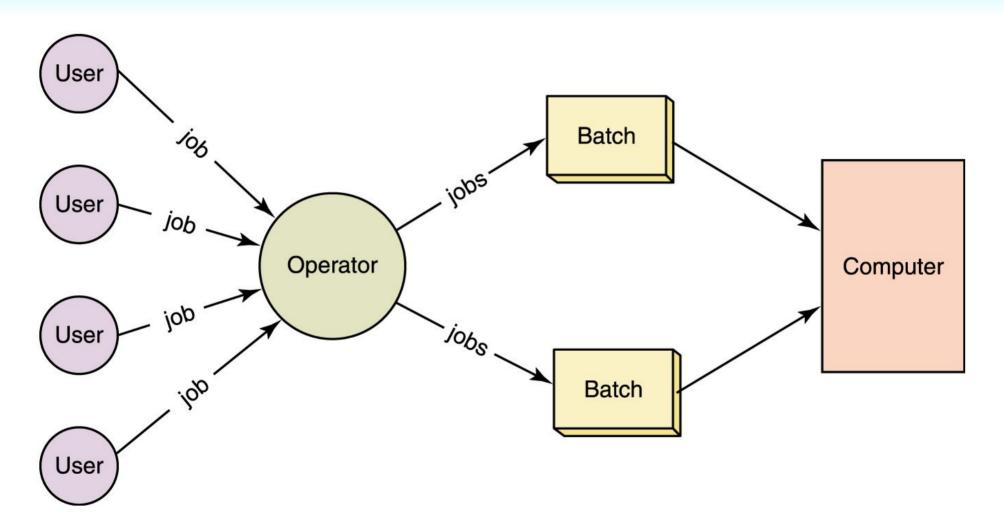


Figure 10.2 In early systems, human operators would organize jobs into batches

Timesharing

- Timesharing system A system that allows multiple users to interact with a computer at the same time
- Multiprogramming A technique that allows multiple processes to be active at once, allowing programmers to interact with the computer system directly, while still sharing its resources
- In a timesharing system, each user has his or her own virtual machine, in which all system resources are (in effect) available for use

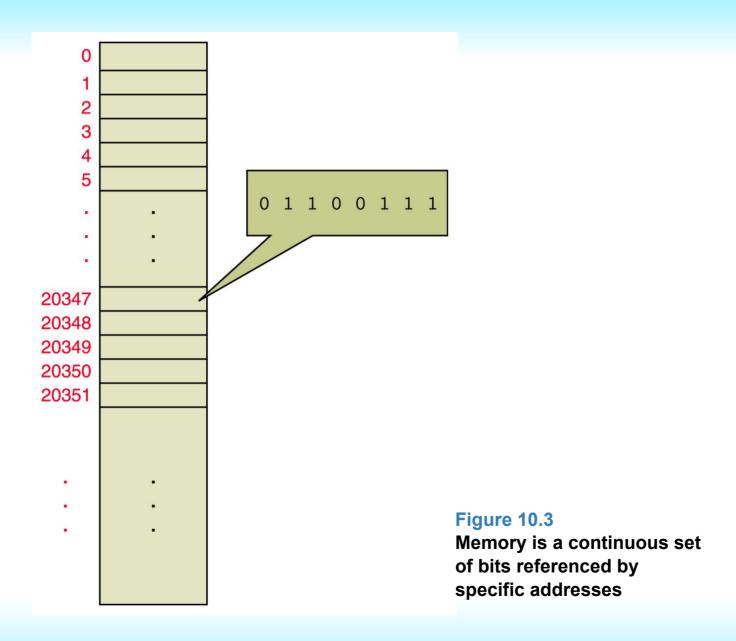
Other Factors

- Real-time System A system in which response time is crucial given the nature of the application
- Response time The time delay between receiving a stimulus and producing a response
- Device driver A small program that "knows" the way a particular device expects to receive and deliver information.

Memory Management

- Operating systems must employ techniques to
 - Track where and how a program resides in memory
 - Convert logical addresses into actual addresses
- Logical address (sometimes called a virtual or relative address) A value that specifies a generic location, relative to the program but not to the reality of main memory
- Physical address An actual address in the main memory device

Memory Management



Single Contiguous Memory Management

Operating system

Application program

- There are only two programs in memory
 - The operating system
 - The application program
- This approach is called single contiguous memory management

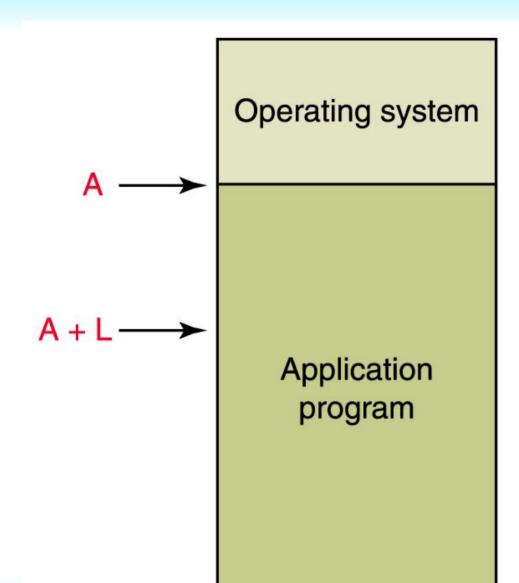
Figure 10.4 Main memory

Main memory divided into two sections

Single Contiguous Memory Management

- A logical address is simply an integer value relative to the starting point of the program
- To produce a physical address, we add a logical address to the starting address of the program in physical main memory

Single Contiguous Memory Management



Logical address L

translates to

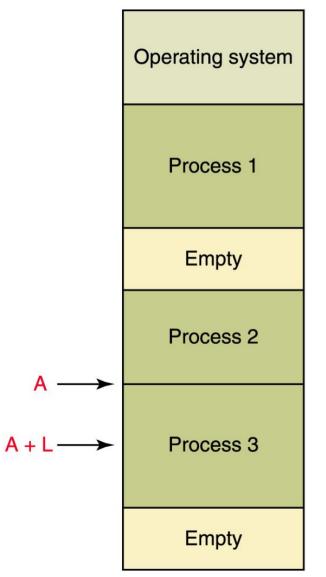
Physical address A + L

Figure 10.5 binding a logical address to a physical one

Partition Memory Management

- Fixed partitions Main memory is divided into a particular number of partitions
- Dynamic partitions Partitions are created to fit the needs of the programs

Partition Memory Management



Base register

Bounds register length

Check: L < length? Yes

Figure 10.6
Address resolution in partition memory management

- At any point in time memory is divided into a set of partitions, some empty and some allocated to programs
- Base register A register that holds the beginning address of the current partition
- Bounds register A register that holds the length of the current partition

Partition Selection Algorithms

Which partition should we allocate to a new program?

- First fit Allocate program to the first partition big enough to hold it
- Best fit Allocated program to the smallest partition big enough to hold it
- Worst fit Allocate program to the largest partition big enough to hold it

- Paged memory technique A memory management technique in which processes are divided into fixed-size pages and stored in memory frames when loaded into memory
 - **Frame** A fixed-size portion of main memory that holds a process page
 - Page A fixed-size portion of a process that is stored into a memory frame
 - Page-map table (PMT) A table used by the operating system to keep track of page/frame relationships

P1 PMT Page Frame 0 5 1 12 2 15 3 7 4 22

P2 PMT	
Page	Frame
0	10
1	18
2	1
3	11

Figure 10.7 A paged memory management approach

•	vierrior y
Frame	Contents
0	
1	P2/Page2
2	
3	
4	
5	P1/Page0
6	
7	P1/Page3
8	
9	
10	P2/Page0
11	P2/Page3
12	P1/Page1
13	
14	
15	P1/Page2

Memory

- To produce a physical address, you first look up the page in the PMT to find the frame number in which it is stored
- Then multiply the frame number by the frame size and add the offset to get the physical address

- Demand paging An important extension of paged memory management
 - Not all parts of a program actually have to be in memory at the same time
 - In demand paging, the pages are brought into memory on demand
- **Page swap** The act of bringing in a page from secondary memory, which often causes another page to be written back to secondary memory

- The demand paging approach gives rise to the idea of **virtual memory**, the illusion that there are no restrictions on the size of a program
- Too much page swapping, however, is called thrashing and can seriously degrade system performance.

Have A Good Night



Autumn Moon by Ansel Adams

